

A note about theory of codes of Goppa

Abstract

In this paper, we study about the theory of field of algebraic functions, a small introduction to theory of codes, highlighting at this point the study of geometric codes of Goppa, and some results involving the definition of floor of a divisor. This paper provides a quota for the minimum distance of a geometric code of Goppa. A such quota involve the floor of a divisor and provide a estimate for the minimum distance of this code.

1 Introduction

We introduce the theory of field of algebraic functions, such as discrete valuation, places, divisors, genus and adeles of a field of algebraic functions, Weil differentials and the Riemann-Roch theorem. We put also the definition of local components of the Weil differentials and the theory of codes. We define the minimum distance of a code and a geometric code of Goppa. Now, we have a motivation for the definition of a geometric code of Goppa; we consider the Reed Solomon code, a vector space over the field finite with q elements \mathbb{F}_q . The geometric codes of Goppa are a generalization of Reed-Solomon codes.

We consider $n = q - 1$ and $\beta \in \mathbb{F}_q$ an element such that

$$\mathbb{F}_q \setminus \{0\} = \{\beta, \beta^2, \dots, \beta^n = 1\}.$$

Keywords: divisors, minimum distance, floor of a divisor, floor quota.

For $k \in \mathbb{Z}$, with $1 \leq k \leq n$, we consider the \mathbb{F}_q -vector space k -dimensional,

$$\mathcal{L}_k := \{f \in \mathbb{F}_q[X] \mid \deg(f) \leq k - 1\},$$

and the application

$$e_v : \mathcal{L}_k \longrightarrow \mathbb{F}_q^n \text{ given by } e_v(f) := (f(\beta), f(\beta^2), \dots, f(\beta^n)) \in \mathbb{F}_q^n.$$

Note that this application is \mathbb{F}_q -linear, since given $\lambda \in \mathbb{F}_q$ and $f, g \in \mathcal{L}_k$, we have that $e_v(f + \lambda g) = e_v(f) + \lambda(e_v(g))$.

Observe that e_v also is injective, because given a polynomial not null $f \in \mathbb{F}_q[X]$ such that $f \in \mathcal{L}_k$, we have that the $\deg(f) \leq k - 1 \leq n - 1 < n$ and hence f has a number of zeros small than n and thus,

$$\text{Ker}(e_v) = \{f \in \mathcal{L}_k \mid e_v(f) = 0\} = \{0\}$$

and so e_v is injective (observe that $\text{Ker}(e_v) = \{0\}$, since if there exists f not null in \mathcal{L}_k such that $e_v(f) = 0$, the polynomial f has n roots $\beta, \beta^2, \dots, \beta^n$, which is absurd). Therefore,

$$C_k := \{e_v(f) \mid f \in \mathcal{L}_k\}$$

is a code of length n and dimension k over the field \mathbb{F}_q . This code is called *Reed-Solomon code*.

Moreover, we speak about the floor of a divisor. We obtain a quota for the minimum distance of a geometric code of Goppa involving the floor of a divisor, which is given in the Theorem 4.6.

2 Preliminary

The contents of this section it follows of [9].

Definition 2.1. We take $0 \neq x \in F$ and we denote by Z , respectively N , the set of zeros, and the set of poles, of x in \mathbb{P}_F . Then we define:

$(x)_0 := \sum_{P \in Z} v_P(x) P$, the divisor of zeros of x ; $(x)_\infty := \sum_{P \in N} -v_P(x) P$, the divisor of poles of x ; $(x) := (x)_0 - (x)_\infty$, the divisor principal of x .

As $v_P(x) > 0$ for all $P \in Z$ it follows that $(x)_0 \geq 0$, and as $v_P(x) < 0$ for all $P \in N$ we have that $-v_P(x) > 0$, and then $(x)_\infty \geq 0$.

Moreover, $(x) = \sum_{P \in \mathbb{P}_F} v_P(x) P$.

Definition 2.2. For a divisor $A \in \mathbb{D}_F$ we define,

$$\mathcal{L}(A) := \{x \in F \mid (x) + A \geq 0\} \cup \{0\}.$$

We have that $\mathcal{L}(A)$, according to Definition 2.2, is a vector space over the field K . The dimension of each K -vector space V will be denoted by $\ell(V)$. Moreover, the vector space $\mathcal{L}(A)$ has dimension finite, for all divisor A of the field of algebraic functions $F|K$.

Lemma 2.3. (see [9]) *We have that:*

- (a) $\mathcal{L}(0) = K$.
- (b) If $A \in \mathbb{D}_F$ and $A < 0$ then $\mathcal{L}(A) = \{0\}$.

Definition 2.4. For $A \in \mathbb{D}_F$, we define the dimension of the divisor A , and we denote by $\ell(A)$, as the dimension of the vector space $\mathcal{L}(A)$, i.e., $\ell(A) = \ell(\mathcal{L}(A))$.

The genus g of the field of algebraic functions $F|K$ is defined by

$$g := \max \{ \deg(A) - \ell(A) + 1 \mid A \in \mathbb{D}_F \}.$$

The genus of $F|K$ is a non-negative integer.

Definition 2.5. An adele of the field of algebraic functions $F|K$ is an application $\alpha : \mathbb{P}_F \rightarrow F$, given by $\alpha(P) = \alpha_P$, such that $\alpha_P \in O_P$, for almost all the places P of $F|K$.

According to the Definition 2.5 we can consider an adele as a sequence in the field F , and therefore we use the notation $\alpha = (\alpha_P)_{P \in \mathbb{P}_F}$ or $\alpha = (\alpha_P)$. The set

$$\mathbb{A}_F := \{ \alpha \mid \alpha \text{ is an adele of } F|K \},$$

is called the **space of adeles** of $F|K$.

This set is a vector space over the field K .

The principal adele of an element $x \in F$ is the adele whose components are all equals to x , and note that $\alpha_P = x \in O_P$, for almost all the places P of $F|K$ make sense, since x not null in F has a finite quantity of poles and zeros. We take the application that take $x \in F$ in the your principal adele; this is a diving of F in \mathbb{A}_F , since is a bijective application, and then we can write $F \hookrightarrow \mathbb{A}_F$.

As we have this diving, the discrete valuation v_P associated to the place P of $F|K$ naturally extends to \mathbb{A}_F , where we put $v_P(\alpha) := v_P(\alpha_P)$, being that we have α_P the P component of the adele α . By definition $v_P(\alpha) \geq 0$, for almost all the places P of $F|K$.

Definition 2.6. For $A \in \mathbb{D}_F$ we define the following set

$$\mathbb{A}_F(A) := \{\alpha \in \mathbb{A}_F \mid v_P(\alpha) \geq -v_P(A), \text{ for all place } P \text{ of } F|K\}.$$

This set is an K -vector subspace of \mathbb{A}_F . We introduce now the concept of Weil differential.

Definition 2.7. A Weil differential of the field of algebraic functions $F|K$ is an application K -linear

$$\omega : \mathbb{A}_F \longrightarrow K,$$

that is null in $\mathbb{A}_F(A) + F$, for some divisor $A \in \mathbb{D}_F$. We define the module of Weil differential of $F|K$ by

$$\Omega_F := \{\omega \mid \omega \text{ is a Weil differential of } F|K\},$$

and we have that Ω_F is an F -vector space.

For $A \in \mathbb{D}_F$ we consider

$$\Omega_F(A) := \{\omega \in \Omega_F \mid \omega \text{ is null in } \mathbb{A}_F(A) + F\}.$$

We have that $\Omega_F(A)$ is a vector space over the field K .

Definition 2.8. For $x \in F$ and $\omega \in \Omega_F$, we define the Weil differential

$$x\omega : \mathbb{A}_F \longrightarrow K \text{ by } (x\omega)(\alpha) := \omega(x\alpha).$$

Definition 2.9. Let P be a place of the field of algebraic functions $F|K$.

- (a) For $x \in F$ we put $\iota_P(x) \in \mathbb{A}_F$ the adele whose P -component is x and all the others components are equals to zero.
- (b) For a Weil differential $\omega \in \Omega_F$ we define the your local component as the function $\omega_P : F \longrightarrow K$ given by

$$\omega_P(x) := \omega(\iota_P(x)).$$

We have that ω_P is an application K -linear.

We want has a divisor for any Weil differential $\omega \neq 0$. We consider for a given ω the set

$$M(\omega) := \{A \in \mathbb{D}_F \mid \omega \text{ is null over } \mathbb{A}_F(A) + F\}.$$

Lemma 2.10. ([9]) *We take $0 \neq \omega \in \Omega_F$. Then, there exists an unique divisor $W \in M(\omega)$ such that $A \leq W$, for any $A \in M(\omega)$.*

The next definition make sense by the Lemma 2.10.

Definition 2.11. (a) The divisor (ω) of a Weil differential ω not null is the unique divisor of the field of algebraic functions $F|K$ such that:

- (1) ω is null over $\mathbb{A}_F((\omega)) + F$.
- (2) If ω is null over $\mathbb{A}_F(A) + F$ then $A \leq (\omega)$.

(b) For $\omega \in \Omega_F \setminus \{0\}$ and P a place of $F|K$, we define $v_P(\omega) := v_P((\omega))$.

(c) A place P is said to be a zero, respectively a pole, of ω if $v_P(\omega) > 0$, respectively $v_P(\omega) < 0$. We say that ω is regular in P if $v_P(\omega) \geq 0$ and ω is called regular, if ω is regular for any place P of $F|K$.

(d) A divisor W is called a canonical divisor of $F|K$ if $W = (\omega)$, for some ω Weil differential not null of $F|K$.

It follows of the Definition 2.11 that:

- (1) $\Omega_F(A) = \{\omega \in \Omega_F \mid \omega = 0 \text{ or } (\omega) \geq A\}$.
- (2) $\Omega_F(0) = \{\omega \in \Omega_F \mid \omega \text{ is regular}\}$.

Theorem 2.12. (see [9])[**Riemann-Roch**] *Let W be a canonical divisor of the field of algebraic functions $F|K$ of genus g . Then, for any divisor A of $F|K$, we have that the dimension of A satisfies the following equality:*

$$\ell(A) = \deg(A) + 1 - g + \ell(W - A).$$

Corollary 2.13. (see [9]) *For a canonical divisor W of the field of algebraic functions $F|K$ of genus g we have that*

$$\deg(W) = 2g - 2 \text{ and } \ell(W) = g.$$

3 On the definition of codes

Let \mathbb{F}_q be a field finite with q elements. We consider the vector space of dimension n over \mathbb{F}_q , \mathbb{F}_q^n , where the elements are of the following form (a_1, \dots, a_n) , with a_i in the field \mathbb{F}_q , for all $i = 1, \dots, n$.

For $a = (a_1, \dots, a_n)$ and $b = (b_1, \dots, b_n)$ in \mathbb{F}_q^n , we take

$$d(a, b) := |\{i \mid a_i \neq b_i\}|.$$

The function d , that we presented above, define a distance in \mathbb{F}_q^n and is called of **Hamming Distance** in \mathbb{F}_q^n . The weight of an element $a \in \mathbb{F}_q^n$ is defined as $\omega(a) := d(a, 0) = |\{i \mid a_i \neq 0\}|$.

Definition 3.1. A code C over the field \mathbb{F}_q is a vector subspace of \mathbb{F}_q^n . The elements of C are called of words of the code. The natural number n , that appear in \mathbb{F}_q^n , is called the length of the code C and $\ell(C)$ is the dimension of C as \mathbb{F}_q -vector space. A code of length n and dimension k is denoted by $[n, k]$. The **minimum distance** of a code $C \neq \{0\}$ is denoted by $d(C)$ and is defined as

$$d(C) := \min \{d(a, b) \mid a, b \in C \text{ and } a \neq b\}.$$

As $d(a, b) = |\{i \mid a_i - b_i \neq 0\}| = d(a - b, 0) = \omega(a - b)$ and $a - b \in C$, since C is a vector space, it follows that the minimum distance is such that

$$d(C) := \min \{\omega(c) \mid 0 \neq c \in C\}.$$

A code $[n, k]$ with minimum distance d will be referred as a code $[n, k, d]$. The canonical inner product over \mathbb{F}_q^n is defined by,

$$\langle a, b \rangle := \sum_{i=1}^n a_i b_i,$$

for $a = (a_1, \dots, a_n)$ and $b = (b_1, \dots, b_n) \in \mathbb{F}_q^n$.

Definition 3.2. If $C \subseteq \mathbb{F}_q^n$ is a code then

$$C^\perp := \{u \in \mathbb{F}_q^n \mid \langle u, c \rangle = 0, \forall c \in C\}$$

is called the code dual of C .

The code C is called self-dual, respectively self-orthogonal, if $C = C^\perp$, respectively, if $C \subseteq C^\perp$. Moreover, we have that $(C^\perp)^\perp = C$.

Now, let $F|\mathbb{F}_q$ be a field of algebraic functions of genus g and let P_1, \dots, P_n be places two to two different of $F|\mathbb{F}_q$ of degree one. Moreover, we take $D = P_1 + \dots + P_n$. And let G be a divisor of $F|\mathbb{F}_q$ such that $\text{supp}(G) \cap \text{supp}(D) = \emptyset$. According to these considerations, the geometric code of Goppa $C_{\mathcal{L}}(D, G)$, associated with the divisors D and G , is defined by:

$$C_{\mathcal{L}}(D, G) := \{(x(P_1), \dots, x(P_n)) \mid x \in \mathcal{L}(G)\} \subseteq \mathbb{F}_q^n.$$

Definition 3.3. The integer $d^* := n - \deg(G)$ is called the **designated distance** of the code $C_{\mathcal{L}}(D, G)$.

We have that the minimum distance d of a geometric code of Goppa is such that $d^* \leq d$.

Now, we define an other code associated with the divisors G and D .

Definition 3.4. Let G and $D = P_1 + \dots + P_n$ be divisors as before (i.e., we have that, for each i , P_i are places two to two different of degree one and $\text{supp}(G) \cap \text{supp}(D) = \emptyset$). Then, we define the code $C_{\Omega}(D, G) \subseteq \mathbb{F}_q^n$ by:

$$C_{\Omega}(D, G) := \{(\omega_{P_1}(1), \dots, \omega_{P_n}(1)) \mid \omega \in \Omega_F(G - D)\}.$$

The following theorem will be used in the next section.

Theorem 3.5. (see [9]) *The codes $C_{\mathcal{L}}(D, G)$ and $C_{\Omega}(D, G)$ are dual, i.e.,*

$$C_{\Omega}(D, G) = C_{\mathcal{L}}(D, G)^{\perp}.$$

4 On the floor of a divisor

In this section we presented the results.

Given a divisor A of a field of algebraic functions $F|K$, there exists an unique divisor of minimum degree B that define the same vector space $\mathcal{L}(A)$, i.e., such that $\mathcal{L}(A) = \mathcal{L}(B)$.

A such divisor is called the floor of the divisor.

4.1 The floor of a divisor

Throughout this section, $F|\mathbb{F}_q$ is a field of algebraic functions on the finite field \mathbb{F}_q which has q elements.

Definition 4.1. Given two divisors A and A' of $F|\mathbb{F}_q$ the greatest common divisor of A and A' is the divisor

$$\text{m.d.c.} (A, A') := \sum_{P \in \mathbb{P}_F} \min \left\{ v_P(A), v_P(A') \right\} P,$$

and the least common multiple between the divisors A and A' is the divisor

$$\text{m.m.c.} (A, A') := \sum_{P \in \mathbb{P}_F} \max \left\{ v_P(A), v_P(A') \right\} P.$$

4.2 The results

Proposition 4.2. *Let G be a divisor of the field of algebraic functions $F|\mathbb{F}_q$ with $\ell(G) > 0$. Suppose that G' is a divisor of $F|\mathbb{F}_q$ of minimum degree such that $\mathcal{L}(G) = \mathcal{L}(G')$. Then, $G \geq G'$. Hence, G' is the unique divisor with this property.*

Proof. By hypothesis, we have that $\mathcal{L}(G) = \mathcal{L}(G') \cap \mathcal{L}(G)$. Thus $\mathcal{L}(G') \cap \mathcal{L}(G) = \mathcal{L}(\text{m.d.c.}(G', G))$: in fact, if $0 \neq x \in \mathcal{L}(G') \cap \mathcal{L}(G)$ then, by Definition 2.2, $(x) + G' \geq 0$, $(x) + G \geq 0$. So, $\min \{v_P((x) + G'), v_P((x) + G)\} = \min \{v_P((x)) + v_P(G'), v_P((x)) + v_P(G)\} \geq 0$, $\forall P \in \mathbb{P}_F$.

On the other hand, it is seen that the minimum value of the set given previous is as it follows, $v_P(x) + \min \{v_P(G'), v_P(G)\}$. Therefore,

$$\sum_P \left(v_P(x) + \min \left\{ v_P(G'), v_P(G) \right\} \right) P \geq 0,$$

for all $P \in \mathbb{P}_F$, i.e., $(x) + \text{m.d.c.}(G', G) \geq 0$. Thus, $x \in \mathcal{L}(\text{m.d.c.}(G', G))$.

Conclusion: $\mathcal{L}(G') \cap \mathcal{L}(G) \subseteq \mathcal{L}(\text{m.d.c.}(G', G))$.

We consider $x \in \mathcal{L}(\text{m.d.c.}(G', G))$. As $\min \{v_P(G), v_P(G')\} \leq v_P(G)$ $\forall P \in \mathbb{P}_F$, it follows that:

$$\sum_P \min \{v_P(G), v_P(G')\} P \leq \sum_P v_P(G) P \Rightarrow 0 \leq (x) + \text{m.d.c.}(G, G') \leq (x) + G \Rightarrow x \in \mathcal{L}(G) = \mathcal{L}(G') \Rightarrow x \in \mathcal{L}(G) \cap \mathcal{L}(G').$$

Conclusion: $\mathcal{L}(\text{m.d.c.}(G', G)) \subseteq \mathcal{L}(G') \cap \mathcal{L}(G)$.

Thus $\mathcal{L}(G) = \mathcal{L}(G) \cap \mathcal{L}(G') = \mathcal{L}(\text{m.d.c.}(G', G))$. Therefore, it follows of the minimum degree of G' that, $\deg(G') \leq \deg(\text{m.d.c.}(G', G))$.

As $\min \{v_P(G'), v_P(G)\} \leq v_P(G')$ for all place P , we obtain

$$\sum_P \min \{v_P(G'), v_P(G)\} P \leq \sum_P v_P(G') P,$$

and then we have that $\text{m.d.c.}(G', G) \leq G'$. Therefore, we have: $\deg(G') \leq \deg(\text{m.d.c.}(G', G))$ and $\text{m.d.c.}(G', G) \leq G'$. Thus, $\deg(\text{m.d.c.}(G', G)) \leq \deg(G')$ and therefore we have that $\deg(G') = \deg(\text{m.d.c.}(G', G))$; as $\text{m.d.c.}(G', G) \leq G'$ we have that, $G' = \text{m.d.c.}(G', G)$. Hence, $G' \leq G$, since $\text{m.d.c.}(G', G) \leq G$.

Suppose now that G' and G'' are two divisors of $F|\mathbb{F}_q$ of minimum degree such that, $\mathcal{L}(G) = \mathcal{L}(G')$ and $\mathcal{L}(G) = \mathcal{L}(G'')$. Now, of G'' divisor of minimum degree such that $\mathcal{L}(G') = \mathcal{L}(G'')$ it follows, of the done previously, that $G' \geq G''$ and of G' divisor of minimum degree such that $\mathcal{L}(G'') = \mathcal{L}(G')$, we have that $G'' \geq G'$. Therefore, $G' = G''$. Hence, there exists an unique divisor G' of $F|\mathbb{F}_q$ of minimum degree such that $\mathcal{L}(G) = \mathcal{L}(G')$. \square

Definition 4.3. Given a divisor G of the field of algebraic functions $F|\mathbb{F}_q$ with $\ell(G) > 0$, the floor of G is the unique divisor G' of $F|\mathbb{F}_q$ of minimum degree such that $\mathcal{L}(G) = \mathcal{L}(G')$. The floor of G will be denoted by $\lfloor G \rfloor$.

4.3 The theorem of the quota of the floor

The next proposition will be used in the sequel.

Proposition 4.4. *If G is a divisor effective of the field of algebraic functions $F|\mathbb{F}_q$ then the floor of G also is a divisor effective. In particular, if G is a divisor effective then the support of $\lfloor G \rfloor$ is a subset of the support of G .*

Proof. As $G \geq 0$, the elements of \mathbb{F}_q belongs to $\mathcal{L}(G)$. Thus, we have that

$$-\min \{v_P(x) \mid x \in \mathcal{L}(G) \setminus \{0\}\} \geq 0.$$

We take $E := \text{m.d.c.}(G + (x) \mid x \in \mathcal{L}(G) \setminus \{0\})$. So, it follows that $\lfloor G \rfloor = G - E$. So, $v_P(\lfloor G \rfloor) = -\min \{v_P(x) \mid x \in \mathcal{L}(G) \setminus \{0\}\}$ the that implies $v_P(\lfloor G \rfloor) \geq 0, \forall P \in \mathbb{P}_F$. Thus, $\lfloor G \rfloor$ is a divisor effective. As $G \geq \lfloor G \rfloor$, we have $v_P(G) \geq v_P(\lfloor G \rfloor), \forall P \in \mathbb{P}_F$, and hence the support of the floor of G is contained in the support of G . \square

Moreover, also will be used in the sequel the following result.

Lemma 4.5. (see [9]) *Let A and B be two divisors of the field of algebraic functions $F|K$ with $A \leq B$. Then, we have that $\mathcal{L}(A) \subseteq \mathcal{L}(B)$ and $\ell(B) - \ell(A) \leq \deg(B) - \deg(A)$.*

Theorem 4.6 (Floor quota). *Let $F|\mathbb{F}_q$ be a field of algebraic functions of genus g . Let $D = P_1 + \dots + P_n$ be a divisor, where P_1, \dots, P_n are places two to two different of $F|\mathbb{F}_q$, where each divisor have degree equal to one, and let $G := H + \lfloor H \rfloor$ be a divisor of $F|\mathbb{F}_q$ such that H is a divisor effective whose support does not contain any of the places P_1, \dots, P_n . Let $E_H := H - \lfloor H \rfloor$ be a divisor. Then, $C_\Omega(D, G)$ is a code $[n, k, d]$ whose minimum distance satisfies:*

$$d = d(C_\Omega(D, G)) \geq \deg(G) - (2g - 2) + \deg(E_H) = 2\deg(H) - (2g - 2).$$

Proof. As H is a divisor effective, we have, according to the Proposition 4.4, that $\lfloor H \rfloor$ is effective and $\text{supp}(\lfloor H \rfloor) \subseteq \text{supp}(H)$. Thus, $G := H + \lfloor H \rfloor$ is a divisor effective and as the supports of the divisors H and D are disjoint, it follows that the supports of the divisors G and D are also disjoint. We chose a Weil differential $\eta \in \Omega_F(G - D)$ such that the code word $c_0 := (\eta_{P_1}(1), \dots, \eta_{P_n}(1))$ is of minimum weight. By definition, the minimum distance of a code is the minimum of the set of the weights of the code words not nulls; thus, we can assume, without loss of generality, that the first d -coordinates of c_0 are not nulls and that the remaining coordinates are nulls. Then, with $D' := P_1 + \dots + P_d$ we obtain, according to Definition 2.11, that $(\eta) \geq G - D'$, since as $\text{supp}(D') \subseteq \text{supp}(D)$ and $(\eta) \geq G - D$ it follows that $v_Q(\eta) \geq v_Q(G - D)$, for all place Q of $F|\mathbb{F}_q$ that belongs to the support of D and then, in particular, this is applied to every place that belongs to the support of D' . Thus, there exists a divisor effective A such that $(\eta) = G - D' + A$ and whose support does not contain any of the places P_1, \dots, P_d . In the equality $(\eta) = G - D' + A$, we take the degree of the divisors on both sides, and hence we have that $\deg((\eta)) = \deg(G) - d + \deg(A)$, and as $\deg((\eta)) = 2g - 2$, (see the Corollary 2.13) since (η) is canonical divisor, it follows that $2g - 2 = \deg(G) - d + \deg(A)$, and therefore $d = \deg(G) - (2g - 2) + \deg(A)$. In order to prove the assertion about the minimum distance quota, it is sufficient to show that $\deg(A) \geq \deg(E_H)$. Observe that $\deg(A) \geq \ell(H + A) - \ell(H) = \ell(H + A) - \ell(\lfloor H \rfloor) \geq \ell(H + A) - \ell(\lfloor H \rfloor + A)$. In fact,
(1) $\deg(A) \geq \ell(H + A) - \ell(H)$, since as $A \geq 0$ we have that $H + A \geq H$ and then, by the Lemma 4.5, it follows that $\ell(H + A) - \ell(H) \leq \deg(H + A) -$

$\deg(H) = \deg(A)$. Therefore, $\deg(A) \geq \ell(H + A) - \ell(H)$.

(2) $\ell(H + A) - \ell(H) = \ell(H + A) - \ell(\lfloor H \rfloor)$, since as $\mathcal{L}(H) = \mathcal{L}(\lfloor H \rfloor)$, we have that $\ell(H) = \ell(\lfloor H \rfloor)$, as desired.

(3) $\ell(H + A) - \ell(\lfloor H \rfloor) \geq \ell(H + A) - \ell(\lfloor H \rfloor + A)$, since as $\lfloor H \rfloor \leq \lfloor H \rfloor + A$ it follows that $\mathcal{L}(\lfloor H \rfloor) \subseteq \mathcal{L}(\lfloor H \rfloor + A)$; thus, $\ell(\lfloor H \rfloor) \leq \ell(\lfloor H \rfloor + A)$ and then $-\ell(\lfloor H \rfloor) \geq -\ell(\lfloor H \rfloor + A)$ and adding $\ell(H + A)$ on both sides of this inequality we have the result.

We show now that $\deg(E_H) = \ell(H + A) - \ell(\lfloor H \rfloor + A)$. We have $W = G - D' + A$ canonical divisor. We take the divisor $H + A$ and the divisor $\lfloor H \rfloor + A$; by the Riemann-Roch Theorem (Theorem 2.12) we obtain the following equalities:

- (1) $\ell(H + A) = \deg(H + A) + 1 - g + \ell(W - H - A)$ and
- (2) $\ell(\lfloor H \rfloor + A) = \deg(\lfloor H \rfloor + A) + 1 - g + \ell(W - \lfloor H \rfloor - A)$.

Thus, $\ell(H + A) - \ell(\lfloor H \rfloor + A) = \deg(H) + \ell(W - H - A) - \deg(\lfloor H \rfloor) - \ell(W - \lfloor H \rfloor - A)$.

Therefore, $\ell(H + A) - \ell(\lfloor H \rfloor + A) = \deg(H - \lfloor H \rfloor) + \ell(W - H - A) - \ell(W - \lfloor H \rfloor - A) = \deg(E_H) + \ell(W - H - A) - \ell(W - \lfloor H \rfloor - A)$.

As $W - A = G - D'$ and $G = H + \lfloor H \rfloor$, we obtain that:

$$\ell(H + A) - \ell(\lfloor H \rfloor + A) = \deg(E_H) + \ell(\lfloor H \rfloor - D') - \ell(H - D').$$

We show now that $\mathcal{L}(\lfloor H \rfloor - D') = \mathcal{L}(H - D')$. As $H - D' \leq H$, it follows that $\mathcal{L}(H - D') \subseteq \mathcal{L}(H)$ and then $\mathcal{L}(H - D') \subseteq \mathcal{L}(\lfloor H \rfloor)$, and thus $\mathcal{L}(H - D') = \mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor)$. Note that $\mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor) = \mathcal{L}(\text{m.d.c.}(H - D'; \lfloor H \rfloor))$.

In fact, for $0 \neq x \in \mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor)$, $(x) + H - D' \geq 0$ and $(x) + \lfloor H \rfloor \geq 0$ and then for all place P we have that $v_P((x) + H - D') \geq 0$ and $v_P((x) + \lfloor H \rfloor) \geq 0$ and thus,

$$\begin{aligned} & \min \{v_P((x) + H - D'); v_P((x) + \lfloor H \rfloor)\} = \\ & v_P((x)) + \min \{v_P(H - D'); v_P(\lfloor H \rfloor)\} \geq 0, \end{aligned}$$

for all place P of $F|\mathbb{F}_q$. Therefore,

$$\sum_P \left(v_P((x)) + \min \{v_P(H - D'); v_P(\lfloor H \rfloor)\} \right) P \geq 0.$$

Thus it follows that, $(x) + \text{m.d.c.}(H - D'; \lfloor H \rfloor) \geq 0$ and then $x \in \mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor))$. Therefore,

$$\mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor) \subseteq \mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor)).$$

If $0 \neq x \in \mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor))$ by definition we have that, $(x) + \text{m.d.c.}(H - D', \lfloor H \rfloor) \geq 0$. We have that, $\min\{v_P(H - D'), v_P(\lfloor H \rfloor)\} \leq v_P(H - D'), \forall$ place P . Thus, $0 \leq (x) + \text{m.d.c.}(H - D', \lfloor H \rfloor) \leq (x) + (H - D')$; and so $x \in \mathcal{L}(H - D') \subseteq \mathcal{L}(\lfloor H \rfloor)$ and then $x \in \mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor)$. Therefore, $\mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor)) \subseteq \mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor)$. So, $\mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor)) = \mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor)$.

By hypothesis, $\text{supp}(H) \cap \text{supp}(D) = \emptyset$ and as D' is a divisor effective such that $D' \leq D$ we have that $\text{supp}(H) \cap \text{supp}(D') = \emptyset$ and as $\text{supp}(\lfloor H \rfloor) \subseteq \text{supp}(H)$ it follows that $\text{supp}(\lfloor H \rfloor) \cap \text{supp}(D') = \emptyset$. By the definition, the maximum common divisor between $H - D'$ and $\lfloor H \rfloor$ is the divisor whose coefficients of the places are the minimum of the coefficients of the places of $H - D'$ and $\lfloor H \rfloor$, and as $H \geq \lfloor H \rfloor$ we conclude that $\text{m.d.c.}(H - D', \lfloor H \rfloor) = \lfloor H \rfloor - D'$.

It follows then that $\mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor)) = \mathcal{L}(\lfloor H \rfloor - D')$ and, as we have $\mathcal{L}(\text{m.d.c.}(H - D', \lfloor H \rfloor)) = \mathcal{L}(H - D') \cap \mathcal{L}(\lfloor H \rfloor) = \mathcal{L}(H - D')$, since $\mathcal{L}(H - D') \subseteq \mathcal{L}(\lfloor H \rfloor)$, we obtain that $\mathcal{L}(\lfloor H \rfloor - D') = \mathcal{L}(H - D')$. Therefore:

$$\begin{aligned} d &= d(C_\Omega(D, G)) \\ &= \deg(G) - (2g - 2) + \deg(A) \\ &\geq \deg(G) - (2g - 2) + \deg(E_H) \\ &= \deg(H) + \deg(\lfloor H \rfloor) - (2g - 2) + \deg(H) - \deg(\lfloor H \rfloor) \\ &= 2\deg(H) - (2g - 2). \end{aligned}$$

□

5 Conclusion

In general, to obtain quotas for the minimum distance of a given code is not a very easy problem. One of the reasons for the interest in Goppa Geometric

codes is that, for this great class of codes, it is possible to evaluate a good quota for the minimum distance.

For more references on this theory, see [1], [2], [3], [4], and [5]. Moreover, see also [6], and, [7].

The author of the article also recommends viewing the following articles [10], [11], and, [12], and, [13].

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